

Data Compression

Why we need data compression?

- Fax machine: 40000 DPSI => 4 million dots per page
- 56 KBPS modem, time to transmit = ?
- Video: 30 pictures per second
- Each picture = 200,000 dots or pixels
- 8-bits to represent each primary color
- Bits required for one picture = ?
- Two hour movie requires = ?

Introduction

- Compression is a way to **reduce the number of bits** in a frame but **retaining its meaning**.
- Decreases space, time to transmit, and cost
- Technique is to identify *redundancy* and to eliminate it
- If a file contains only capital letters, we may encode all the 26 alphabets using 5-bit numbers instead of 8-bit ASCII code

Introduction

- If the file had n -characters, then the **savings** = $(8n-5n)/8n \Rightarrow 37.5\%$

Frequency Dependent Codes

- Not all the characters appear with same frequency, **some are more prevalent** than the others
- **Frequently appearing** characters could be assigned **shorter** codes than the others => results in reduced number of bits
- Such codes are examples of frequency dependent code

Frequency Dependent Codes

- **Huffman code:** (illustrated with a manageable example)

Letter	Frequency (%)
A	25
B	15
C	10
D	20
E	30

Frequency Dependent Codes

- **Huffman code:** Code formation
 - **Assign** weights to each character
 - **Merge** two lightest weights into one root node with sum of weights (if multiple? Not unique code)
 - **Repeat** until one tree is left
 - **Traverse** the tree from root to the leaf (for each node, assign 0 to the left, 1 to the right)

Frequency Dependent Codes

- **Huffman code:** Code Interpretation
 - **No prefix property:** code for any character never appears as the prefix of another code (Verify)
 - Receiver continues to receive bits until it finds a code and forms the character
 - 01110001110110110111 (**extract** the string)

Frequency Dependent Codes

- **Arithmetic compression:** is based on Interpreting a **character-string** as a single real number

Letter	Frequency (%)	Subinterval [p, q]
A	25	[0, 0.25]
B	15	[0.25, 0.40]
C	10	[0.40, 0.50]
D	20	?
E	30	?

Frequency Dependent Codes

- **Arithmetic compression:** Coding 'CABAC'
- **Generate** subintervals of decreasing length, subintervals depend uniquely on the string's characters and their frequencies.
- Interval $[x, y]$ has **width** $w = y - x$, the **new interval** based on $[p, q]$ is $x = x + w.p$, $y = x + w.q$
- **Step 1:** 'C' 0.....0.4.....0.5.....1
based on $p = 0.4$, $q = 0.5$

Frequency Dependent Codes

- Step 2: 'A' 0.4.....0.425.....0.5

based on $p = 0.0$, $q = 0.25$

- Step 3: 'B'

0.4.....0.40625.....0.41.....0.425

based on $p = 0.25$, $q = 0.4$

Step 4: 'A'

Step 5: 'C'

...0.406625... 0.4067187...

Final representation (midpoint)?

Frequency Dependent Codes

- **Arithmetic compression:** Extracting 'CABAC'

N	Interval[p, q]	Width	Character	N-p	(N-p)/width
0.4067	0.4 – 0.5	0.1	C	0.0067	0.067
0.067	0 – 0.25	0.25	A	0.067	0.268
0.268	0.25 – 0.4	0.15	B	0.018	0.12

?

?

When to stop? A **terminal character** is added to the original character set and encoded. During decompression, once it is encountered the process **stops**.

Run Length Encoding

- Huffman code **requires**:
 - frequency values
 - bits are grouped into characters or units

Many items **do not fall** into such category

- machine code files
- facsimile Data (bits corresponding to light or dark area of a page)
- video signals

Run Length Encoding

- For such files, **RLE** is used.
- 'Instead of sending long runs of '0's or '1's, it sends only **how many** are in the run.'
- 70%-80% space is white on a typed character space, so RLE is useful.

Run Length Encoding

- Runs of the same bit
- In facsimile Data, there are many '0's (white spots) → transmit the run-length as fixed size binary integer
- Receiver generates proper number of bits in the run and inserts the *other* bit in between
- 14 zeros, **1**, 9 zeros, **11**, 20 zeros, **1**, 30 zeros, **11**, 11 zeros (number of zeros encoded in 4-bits)

Run Length Encoding

- Runs of the same bit
- Code: 1110 1001 0000 1111 0101 1111 1111
0000 0000 1011
- (next **value after 1111** is added to the run)
- SAVINGS IN BITS: ?
- If the stream started with '1' **instead**?
- Best when there are many long runs of zeros, with increased frequency of '1's, becomes less efficient.

Run Length Encoding

- Runs with different characters
- Send the actual character with the run-length
- HHHHHHHHUFFFFFFFFFYYYYYYYYYYDGGGGG
- code = 7, H, 1, U, 9, F, 11, Y, 1, D, 5, G
- SAVINGS IN BITS (considering ASCII): ?

Run Length Encoding

- Facsimile Compression
- ITU standard (A4 document, 210 by 297 mm)
- 1728 pixels per line
- If 1 bit for each pixel, then over 3 million bits for each page
- A typical page contains many consecutive white or black pixels -> RLE

Run Length Encoding

- Run lengths may vary from 0 to 1728 -> **many Possibilities** and inefficiency with a fixed size code
- Some runs occur more frequently than others, e.g. most typed pages contain 80% white pixels, spacing between letters is fairly consistent
- => probabilities of certain runs are predictable
- => Frequency dependent code on run lengths

Run Length Encoding

- Some Facsimile compression codes
(Terminating, less than 64)

Pixels in the run	Code: White	Code: Black
0	00110101	0000110111
1	000111	010
2	0111	11
3	1000	10
10	00111	0000100
20	0001000	00001101000

Run Length Encoding

- Some Facsimile compression codes (**Make up, greater than or equal to 64**)

Pixels in the run	Code: White	Code: Black
64	11011	0000001111
128	10010	000011001000
256		
512		

129 white: Savings:

No-prefix property, better compression for long-runs

Relative Encoding

- Relative Encoding:
- Some applications may not benefit from the above: video image -> little repetitive within, but much repetition from one image to the next
- Differential encoding is based on coding only the **difference** from one to the next

Relative Encoding

- Relative Encoding:

• 1 2 3 4	1 3 3 4	0 1 0 0
2 5 3 7	2 5 3 7	0 0 0 0
3 6 4 8	3 6 4 7	0 0 0 -1
4 7 5 9	3 7 5 9	-1 0 0 0
1 st Frame	2 nd Frame	Difference

Resulting **difference** can be RLE.

Image Representation

- **BW** pixels each represented by 8-bit level
 - Color composed of R, G, B primaries, each is represented by 8-bit level
- > Each color pixel can be represented by one of $2^8 \cdot 2^8 \cdot 2^8 = 2^{24}$ colors

VGA screen: 640 * 480 pixels

-> $640 * 480 * 24 = 7,372,800$ bits

Image Compression

- JPEG compression – both for grayscale and color images
- Previous compression methods were **lossless** – it was possible to recover all the information from the compressed code
- JPEG is **lossy**: image recovered may not be the same as the original

JPEG Compression

- It consists of three phases: Discrete Cosine Transform (DCT), Quantization, Encoding.
- **DCT:**
Image is divided into blocks of 8×8 pixels
For grey-scale images, pixel is represented by 8-bits
For color images, pixel is represented by 24-bits or three 8-bit groups

JPEG Compression

- DCT takes an 8*8 matrix and produces another 8*8 matrix.

- $$T[i][j] = 0.25 C(i) C(j) \sum_x \sum_y P[x][y] \cos \left((2x+1)i\pi/16 \right) * \cos \left((2y+1)j\pi/16 \right)$$

$$i = 0, 1, \dots, 7, j = 0, 1, \dots, 7$$

$$C(i) = 1/\sqrt{2}, i = 0$$

$$= 1 \text{ otherwise}$$

T contains values called **'Spatial frequencies'**

JPEG Compression

- Spatial frequencies directly relate to *how much the pixel values change* as a function of their positions in the block
- $T[0][0]$ is called the **DC coefficient**, related to average values in the array/matrix, $\text{Cos } 0 = 1$
- Other values of T are called **AC coefficients**, cosine functions of higher frequencies

JPEG Compression

- Case 1: all **P's are same** => image of single color with no variation at all, AC coefficients are all zeros.
- Case 2: **little variation in P's** => many, not all, AC coefficients are zeros.
- Case 3: **large variation in P's** => a few AC coefficients are zeros.

JPEG Compression

P-matrix

T-matrix

- | | | | | | | | | | | | |
|-----|----|----|----|-----|-----|------|------|---|-----|---|-----|
| 20 | 30 | 40 | 50 | 60 | ... | 720 | -182 | 0 | -19 | 0 | ... |
| 30 | 40 | 50 | 60 | 70 | | -182 | 0 | 0 | 0 | 0 | |
| 40 | 50 | 60 | 70 | 80 | | 0 | 0 | 0 | 0 | 0 | |
| 50 | 60 | 70 | 80 | 90 | | -19 | 0 | 0 | 0 | 0 | |
| 60 | 70 | 80 | 90 | 100 | | 0 | 0 | 0 | 0 | 0 | |
| ... | | | | | | ... | | | | | |

'Uniform color change and little fine detail,
easier to compress after DCT'

JPEG Compression

P-matrix

- 100 150 50 100 100 ...
200 10 110 20 200
10 200 130 30 200
100 10 90 190 120
10 200 200 120 90
....

T-matrix

835 15 -17 59 5...
46 -60 -36 11 14
-32 -9 130 105 -37
59 -3 27 -12 30
50 -71 -24 -56 -40
....

'Large color change, difficult to compress after DCT'

JPEG Compression

- To **restore**, inverse DCT (IDCT) is performed:

- $$P[x][y] = 0.25 \sum \sum C(i) C(j) T[i][j] \text{Cos} \\ (2x+1)i\pi/16 * \text{Cos} (2y+1)j\pi/16$$

$$x = 0, 1, \dots, 7, y = 0, 1, \dots, 7$$

- # Can write a C-program to apply DCT on a P-array (8*8) to obtain T-array and also IDCT on T-array to recover P-array.

JPEG Compression

- **Quantization:** Provides an way of **ignoring** small differences in an image that may not be perceptible.
- Another array Q is obtained by dividing each element of T by some number and rounding-off to nearest integer => **loss**

JPEG Compression

T-matrix

- 152 0 -48 0 -8...
0 0 0 0 0
-48 0 38 0 -3
0 0 0 0 0
-8 0 -3 0 13
...

Q-matrix

- 15 0 -5 0 -1...
0 0 0 0 0
-5 0 4 0 0
0 0 0 0 0
-1 0 0 0 1
...

'Divide by 10 and round-off' => 'creates fewer distinct numbers and more consistent pattern'

JPEG Compression

- Can we **recover** by multiplying the elements with 10?
- If the loss is **minimal**, the vision system may not notice.
- Dividing T-elements by the **same number** is not practical, may result in **too much** loss.
- **Retain** the effects of **lower spatial** frequencies as much as possible – less subtle features noticed if changed

JPEG Compression

- **Upper left** elements to be divided with **smaller** values
- **Lower right** elements – higher spatial frequencies, finer details - to be divided with **larger** values
- To define a quantization array U, then
$$Q[i][j] = \text{Round} (T[i][j] \div U[i][j]), i = 0, 1, \dots, 7, j = 0, 1, \dots, 7$$

JPEG Compression

- $U = \begin{matrix} 1 & 3 & 5 & 7 & 9 & \dots \\ 3 & 5 & 7 & 9 & 11 & \\ 5 & 7 & 9 & 11 & 13 & \\ 7 & 9 & 11 & 13 & 15 & \\ 9 & 11 & 13 & 15 & 17 & \\ \dots & & & & & \end{matrix}$ $Q = \begin{matrix} 152 & 0 & -10 & 0 & -1 & \dots \\ 0 & 0 & 0 & 0 & 0 & \\ -10 & 0 & 4 & 0 & 0 & \\ 0 & 0 & 0 & 0 & 0 & \\ -1 & 0 & 0 & 0 & 0 & \\ \dots & & & & & \end{matrix}$

'Q can be well-compressed due to redundant elements'

JPEG Compression

- Encoding 152 0 -10 0 -1 ...

- linearize two 0 0 0 0 0

dimensional data -10 0 4 0 0

and compress it 0 0 0 0 0

-1 0 0 0 0

...

- Row-wise (shorter runs)
- Zigzag (**longer runs**, higher spatial frequencies are gathered together)

JPEG Compression

- Many such 8×8 arrays, adjacent blocks with little difference => **more potential** for compression
- JPEG may provide 95% compression (depends on image and quantization array)
- GIF (Graphic Image Format) reduces color to 256. Best suited for a **few colors** and **sharp boundaries** (charts, lines). Not good for variations and shading – full color photo

Multimedia Compression

- JPEG compresses **still pictures**, motion pictures are compressed by MPEG.
 - Still picture contains 7, 372,800 bits.
 - Compression ratio **20:1** reduces the bits to 368,640
 - With 30 images per second, bits to be handled 11,059,200.
- => **Huge** data transmission if channel is shared

Multimedia Compression

- Based on JPEG compression and **relative encoding**
- Usually **little difference** from one to the next frame => suitable for compression
- A **completely new** scene can not be compressed this way. Also, not suited for moving objects unleashing other objects.
- Three frames: I, P, B.

Multimedia Compression

- **I** (intra-picture) - frame: Just a JPEG encoded image.
- **P** (predicted) – frame: Encoded by computing the differences between a current and a previous frame.
- **B** (bidirectional) - frame: Similar to P-frame except that it is interpolated between previous and future frame.

Multimedia Compression

- I-frames must appear **periodically** in any frame sequence, otherwise:
 - any error introduced in any frame is propagated in all subsequent frames if relative differences are sent
 - in broadcast applications, if one tunes-in late, then he has nothing to compare to

Multimedia Compression

- Typical sequence:

order in which to transmit: I – P – I

P is sandwiched between groups of B

P is the difference from the previous I-frame

B is interpolated from the nearest I and P

order in which to display I – B – B – P – B – B – I

Multimedia Compression

- Coding P:
- Motion compensated Prediction for P-frames
- Divide the image into macro-blocks of 256 pixels (16×16)
- Chrominance arrays are reduced to groups of 4 pixels (becomes 8×8 matrix) representing average values => **loss**, not perceptible

Multimedia Compression

- An algorithm examines each macro-block of P-frame and locates a **best matching** macro-block in the prior I-frame
- Not necessarily in the same relative position (in the **vicinity**)
- Once located, the algorithm calculates the difference and **motion vector** (displacement)
- The above information is encoded and transmitted

Multimedia Compression

- At the decoding end, the **motion-vector** is used to determine the **position** and the *difference* is used to reconstruct the macro-block

Multimedia Compression

- **MP3** (MPEG Level 3) – compression protocol for audio
- Audible range: 20 – 20 KHz
- PCM uses 16 bit samples @ 44.1 KHz
- 1-sec of PCM audio needs: $16 * 44.1 * 1000$ bits
- For two channel stereo: Twice
- For 1-minute audio: 60 times -> 88 Mbits

Multimedia Compression

- MP3 is based on a 'psychoacoustic model' – what we can **hear** and what we can **distinguish**
- '**Auditory masking**' – capture an audio signal, determine what we can not hear, remove those components from the stream, digitize what is left
- '**Sub-band coding**' – decompose the original signal into non-overlapping frequency ranges, create bit-stream for each band

Multimedia Compression

- Sub-bands are encoded **differently**:
- Sub-bands with **loud** signals need good resolution, are coded with **more bits**
- Near-by sub-bands with **weak** signals are effectively masked by louder signals, need less resolution, are coded with **fewer** bits, that results in ***compression***